

Probabilistic Validation of Computer System Security

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(Joint work with DPASA Project Team)

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Everyone says it is important, few approaches exist ...

- Security metrics were an important problem in the 2005 INFOSEC Research Council Hard Problems List
- New security metrics that are linked to the business were ranked first among six key security imperatives developed by over twenty Fortune 500 firms
- New regulatory requirements of Sarbanes-Oxley and the Basel II Accord have created more urgency for metrics that integrate security risk with overall business risk
- Almost every critical infrastructure roadmap lists security metrics as a critical challenge
- The list goes on ...



March 21, 2007 - ReSIST Open Workshop, Budapest, Hungary

Security Validation Truths ...

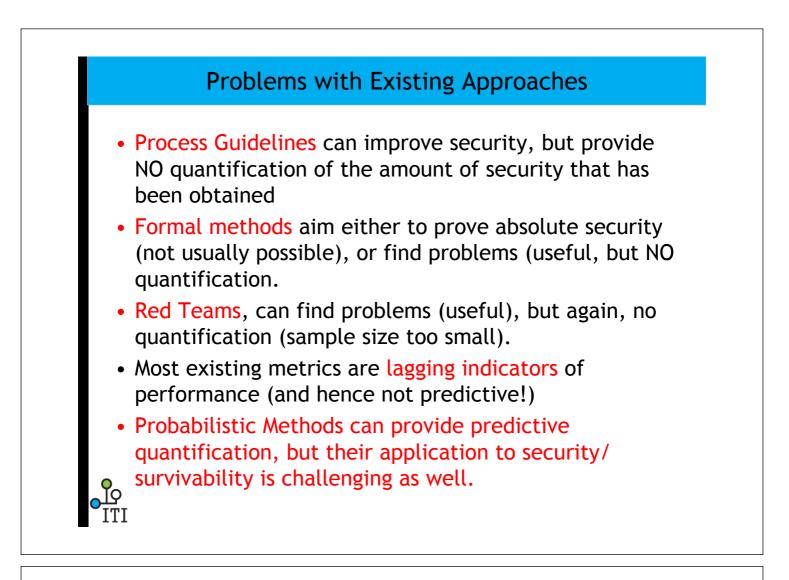
- Security is no longer absolute
- Trustworthy computer systems/networks must operated through attacks, providing proper service in spite of possible partially successful attacks
- Intrusion tolerance claims to provide this ability
- If security is not absolute, quantification of the "amount" of security that a particular approach provides is essential
- Quantification can be useful in:
 - A *relative* sense, to choose amount alternate design alternatives
 - In an *absolute* sense, to provide guarantees to users



Existing Security Validation Approaches

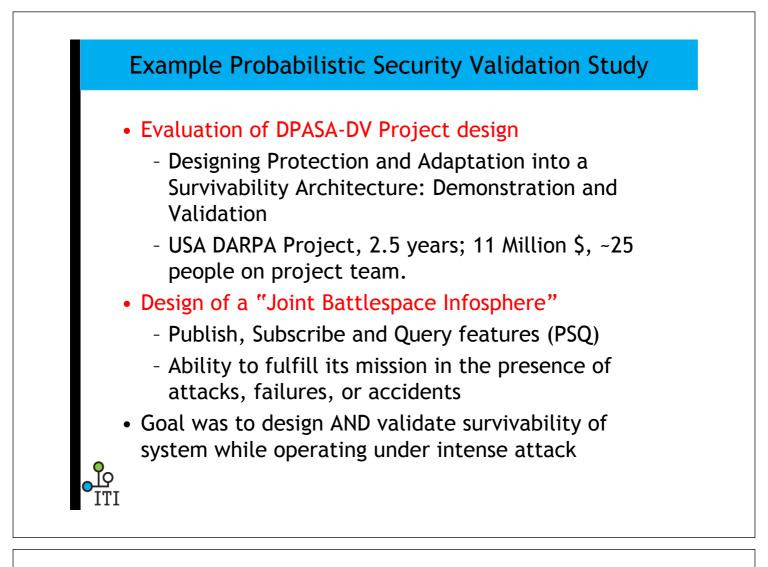
- Most traditional approaches to security validation have focused on and specifying procedures that should be followed during the design of a system (e.g., the Security Evaluation Criteria [DOD85, ISO99]).
- When quantitative methods have been used, they have typically either been based on:
 - formal methods (e.g., [Lan81]), aiming to prove that certain security properties hold given a specified set of assumptions, or
 - been quite informal, using a team of experts (often called a "red team," e.g. [Low01]) to try to compromise a system.

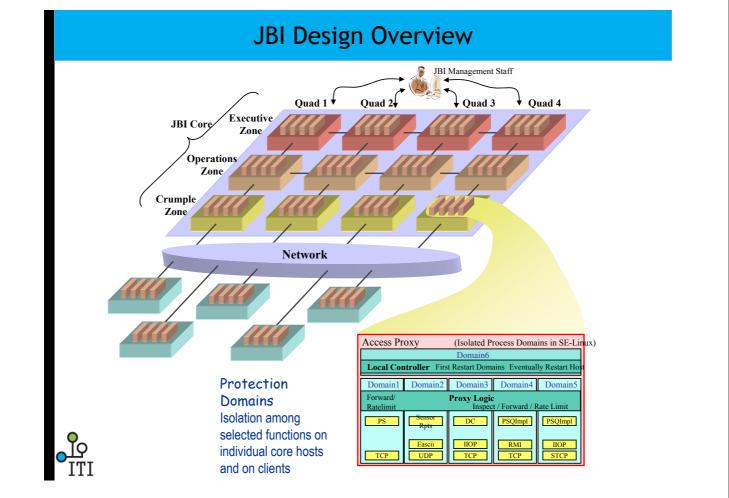


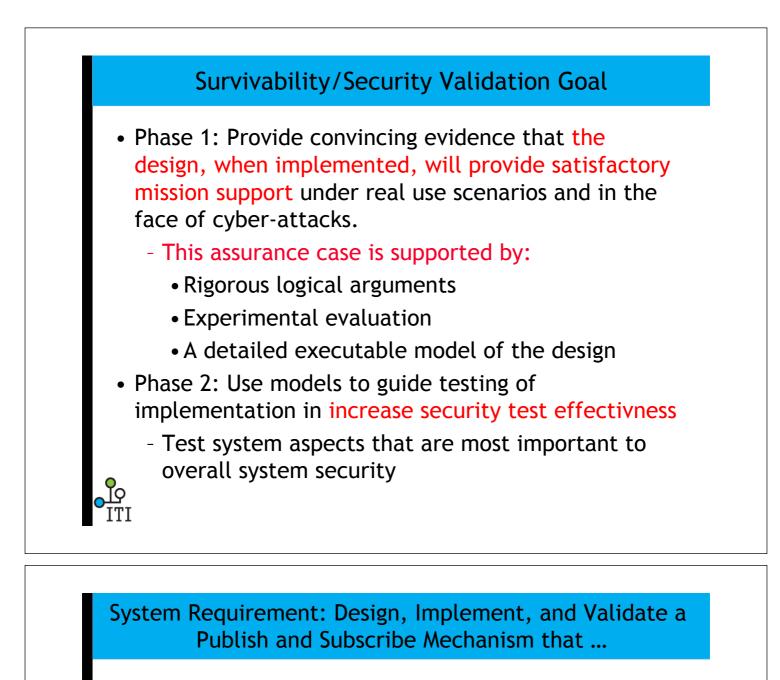


Security Quantification Challenges

- How can the behavior of attackers be quantified?
 - How accurately does this need to be done?
 - At what level of detail?
- How should security/survivability measures be specified?
 - Are new measures needed?
- If relative measures are desired, can they be shown to be robust across a wide variety of situations?
 - Robustness is key to good design
- How accurately can absolute measures be estimated?
- Can quantification aid in security testing?
 - Knowing where to focus testing is key
- Can a notion of "coverage" be developed?
 - If so, testing can produce quantitative results

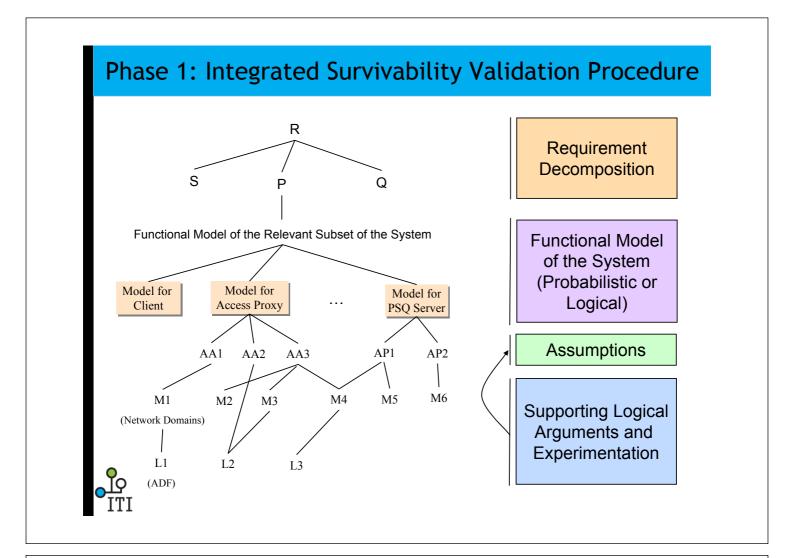


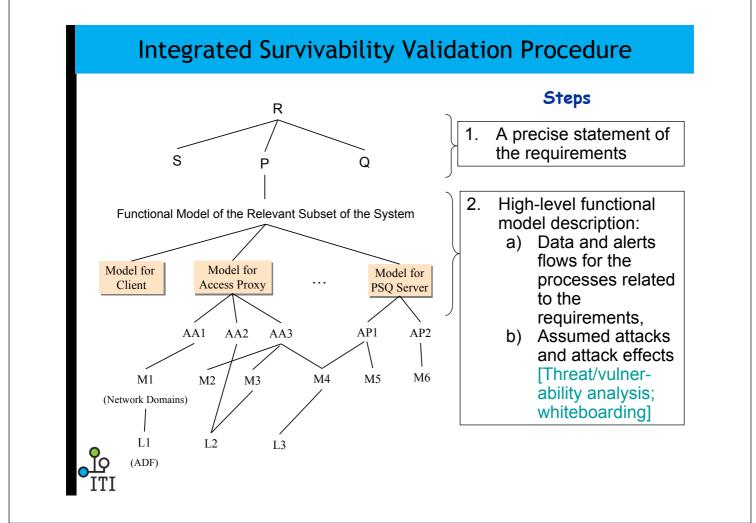


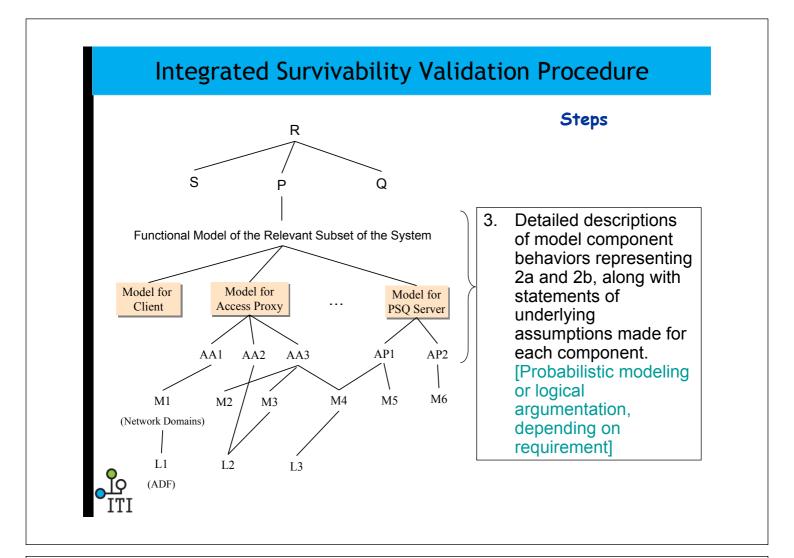


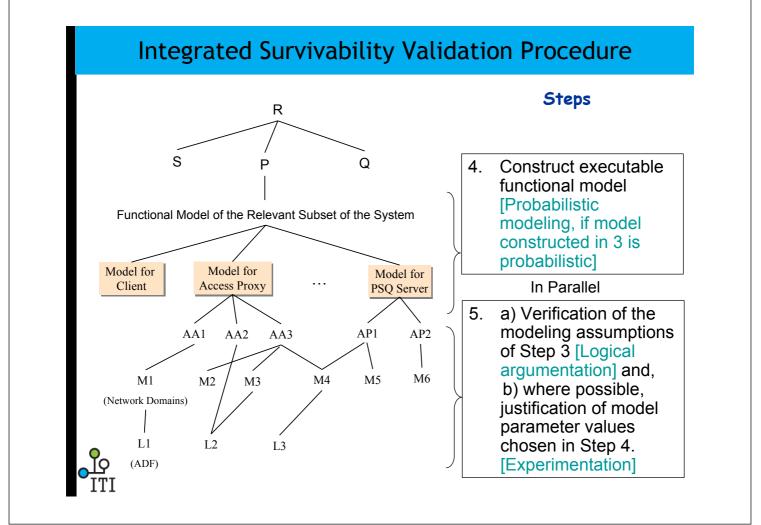
- Provides 100% of critical functionality when under sustained attack by a "Class-A" red team with 3 months of planning
- Detects 95% of large scale attacks within 10 mins. of attack initiation and 99% of attacks within 4 hours with less than 1% false alarm rate
- Displays meaningful attack state alarms. Prevent 95% of attacks from achieving attacker objectives for 12 hours
- Reduces low-level alerts by a factor of 1000 and display meaningful attack state alarms.

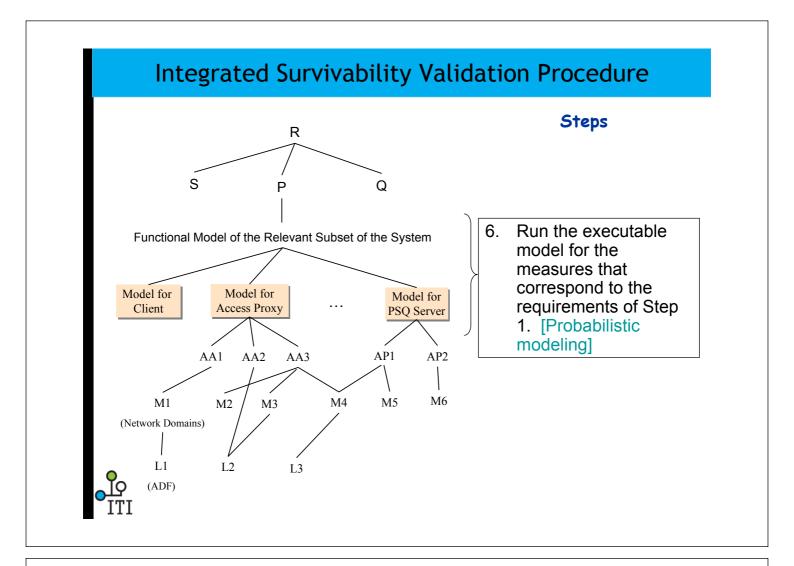
Shows survivability versus cost/performance tradeoffs

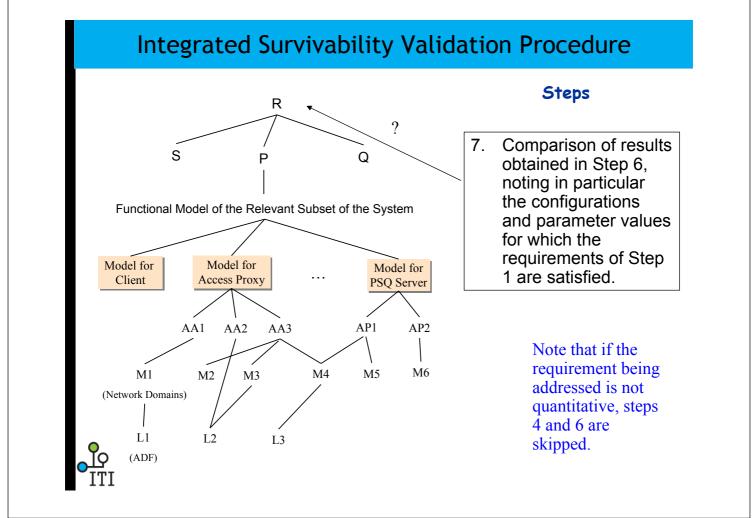


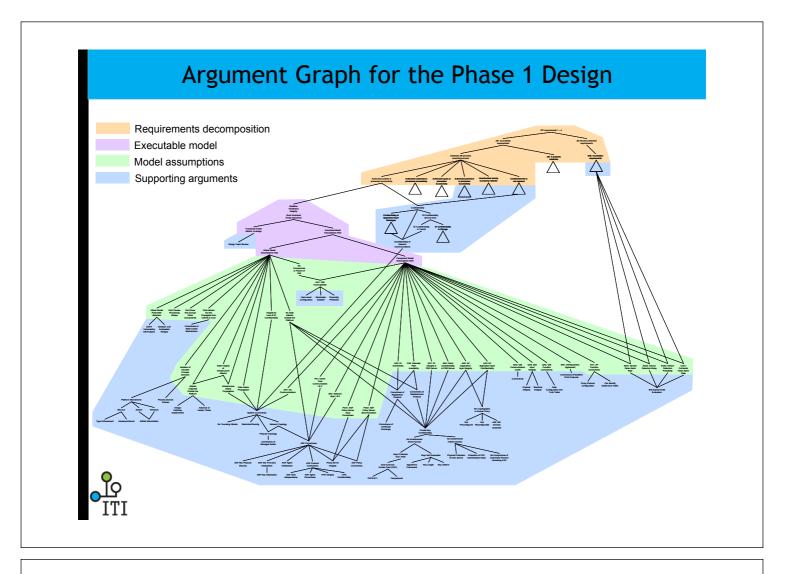






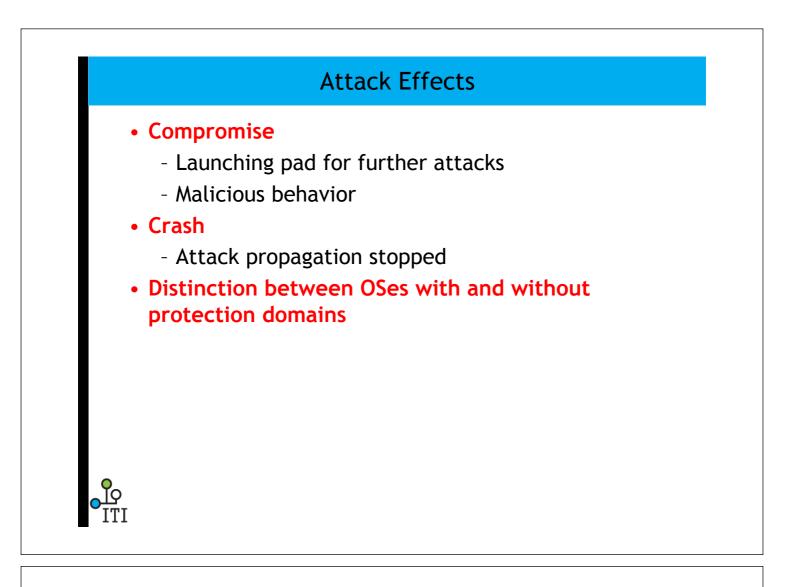






Attack Model Description

- Consider effects of attacks, not attacks themselves
- Attack propagation
 - MTTD: mean time to **discovery** of a vulnerability
 - MTTE: mean time to exploitation of a vulnerability
- 3 types of vulnerabilities:
 - Infrastructure-Level Vulnerabilities → attacks in depth
 - OS vulnerability
 - Non-JBI-specific application-level vulnerability
 - p_{common} : common-mode failure
 - Data-Level Vulnerabilities → attacks in breadth
 - Using the application data of JBI software
 - Across process domains
 - flaw in protection domains



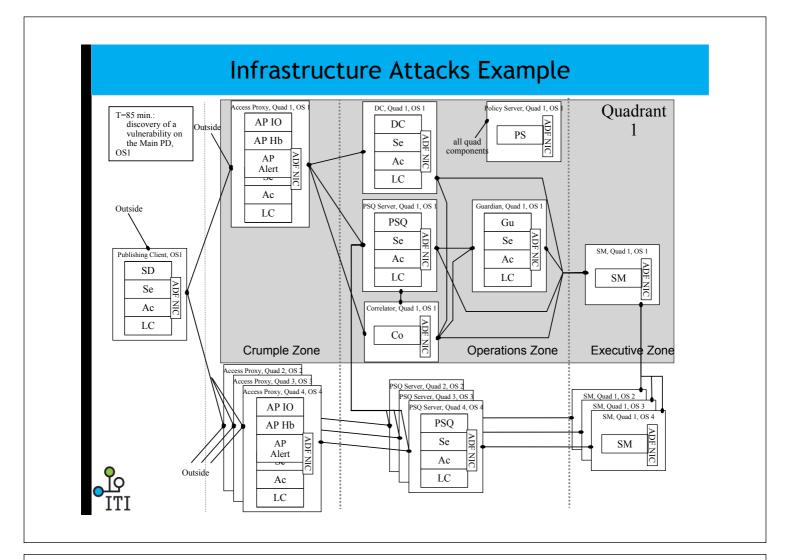
Attack Response

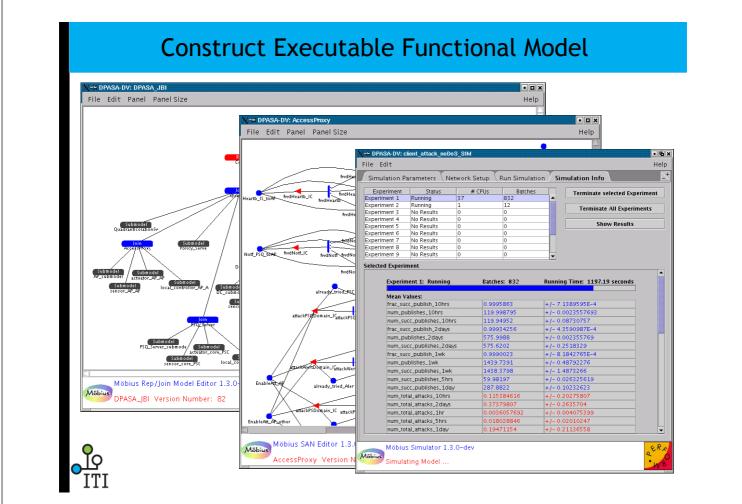
- Intrusion Detection
 - p_{detect} =0 if the sensors are compromised
 - p_{detect} > 0 otherwise.

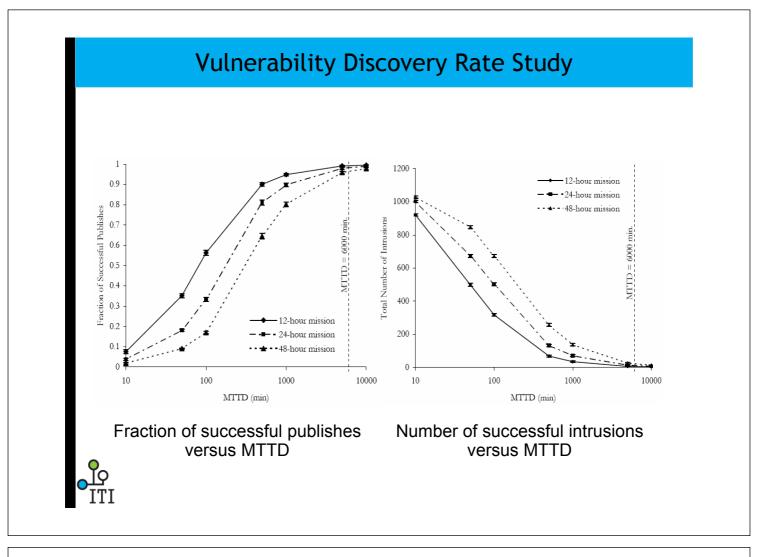
Attack Responses

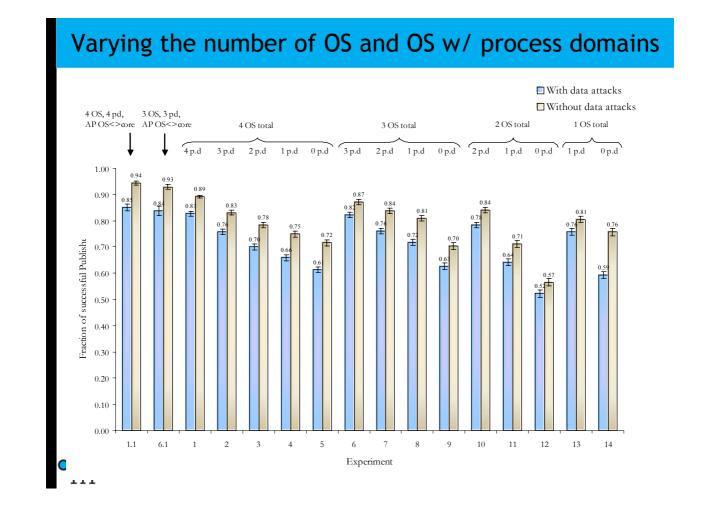
- Restart Processes
- Secure Reboot
- Permanent Isolation











Phase 2: Improving (and Validating) the Implementation

Objectives:

- Improve the system's survivability
- Conduct specific system-level validation tasks
- Address all of the system-level concepts and mechanisms that may contribute to improvement, e.g., protocols and application scenarios

Main Idea:

- Think like an attacker
 - Examine whether a given attacker goal can be achieved
 - If so, alter the implementation so as to preclude such achievement

Procedure:

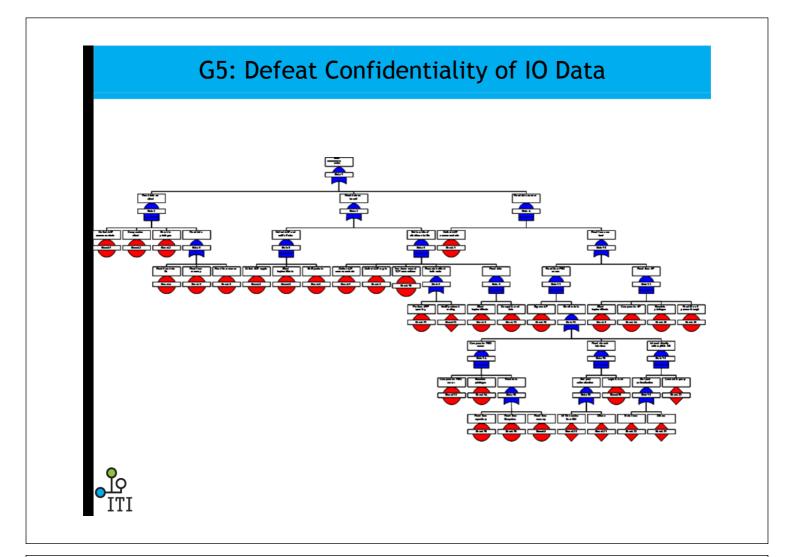
TTT

- Top-down, beginning with a specific high-level attacker goal
- Critical steps of the high-level attack tree are elaborated further
- as sub-trees, down to a level that admits adversarial testing.

Attacker Goals

- We considered the following attacker goals:
 - G1: Prevent client publish
 - G2: Prevent IO delivery to client (Subscription)
 - G3: Prevent a successful query operation
 - G4: Prevent a successful client registration
 - G5: Defeat confidentiality of IO data
 - G6: Modify IO data
 - G7: Modify data in repository





G5: Attack Steps/Minimal Attacks

BASIC	Defeat ADF access control	
	Deleal ADF access control	7,8,9
BASIC	Compromise client	5,3,2,1
UNDEVELOPED	Escalate privilege	4,3,2,1
BASIC	Read from data file	6,3,2,1
BASIC	Read from memory	16 , 21 , 19 , 1
BASIC	Read from screen	16 , 20 , 19 , 1
BASIC	Defeat ADF crypto	16 , 21 , 22 , 1
BASIC	Steal key/certificate	16 , 23 , 22 , 1
BASIC	Sniff packets	
UNDEVELOPED	Tear down current TCP connections	
BASIC	Perform ARP spoofing	
UNDEVELOPED	Modify network routing	
BASIC	Decrypt & read data	
BASIC	Compromised PSQ server	
BASIC	Bypass AP	
BASIC	Read from filesystem	
BASIC	Read from repository	
BASIC	Login & read	
UNDEVELOPED	MITM session from SM	
UNDEVELOPED	Others	
UNDEVELOPED	Connect & query	
UNDEVELOPED	Brute force	
BASIC	Compromise AP	
BASIC	Read IO as it passes through	
	BASIC BASIC BASIC BASIC BASIC BASIC UNDEVELOPED BASIC BASIC BASIC BASIC BASIC BASIC UNDEVELOPED UNDEVELOPED UNDEVELOPED UNDEVELOPED BASIC	BASICRead from data fileBASICRead from memoryBASICRead from screenBASICDefeat ADF cryptoBASICSteal key/certificateBASICSniff packetsUNDEVELOPEDTear down current TCP connectionsBASICPerform ARP spoofingUNDEVELOPEDModify network routingBASICDecrypt & read dataBASICDecrypt & read dataBASICBypass APBASICRead from filesystemBASICRead from repositoryBASICLogin & readUNDEVELOPEDMITM session from SMUNDEVELOPEDOthersUNDEVELOPEDBrute forceBASICCompromise AP

